

## → Digital Signature: -

Message authentication protects two parties who exchange the message from an third party.

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however it doesn't protect each other to solve the dispute between them. The digital signature requirements are: -

- 1- it must be a bit pattern that depend on the message being signed.
- 2- the signature must use some unique information to the sender to prevent both forgery <sup>تکرار</sup> & denial <sup>انکار</sup>.

~~various~~  
various types of digital signature

1- Direct digital signature: -

→ in this method the sender attaches the signature with the msg. the receiver can verify the

# digital signature for authentication purpose

2. arbitrated digital signature.  
the problem with direct D.S can be solved with the help of the third party known as arbitrator.

there are 3 techniques in generation arbitrated signature.

1- conventional cryptography, arbitrator can see the msg.

$$X \rightarrow A : M \parallel E_{K_{XA}} [ID_X \parallel H(M)]$$
$$A \rightarrow Y : E_{K_{AY}} [ID_X \parallel M \parallel E_{K_{XA}} [ID_X \parallel H(M)]]$$

miss 

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 X : generates the msg along with the signature, but the signature is encrypted by using shared key between X & A.

the arbiter will decrypt the signature & generates anew msg which is encrypted by using shared key between A & Y.

→ the receiver y will accept the msg without verifying the signature because the msg has been received from trusted party. arbiter.

~~→ in this case~~

In case of dispute arbiter will resolve the problem by using signature.

(2) Conventional cryptography arbiter doesn't

see message:

$$x \rightarrow A : ID_x : E_{k_{xy}}[M] : E_{k_{xa}}[ID_x : H(E_{xy}(M))]$$

$$A \rightarrow Y : E_{ky}[ID_x : E_{k_{xy}}(M) : E_{k_{xA}}[ID_x : H(E_{k_{xy}}(M))]]$$

3) public key encryption arbitrator  
doesn't see the msg:-

$x \rightarrow A: ID_x, E_{K_{R_x}}[ID_x, E_{K_{y_j}}[$   
 $E_{K_{R_x}}(M)]]$

$A \rightarrow Y: E_{K_{R_A}}[ID_x, E_{K_{y_j}}(E_{K_{R_x}}(M))]$

here x will encrypt the msg by using  
his private key and again encrypt  
with public key of y.

thus both authentication and privacy  
can be achieved and along with  
identity encryption with private key  
of x and sends the message to  
arbitrator.

→ arbitrator will decrypt the message  
with public key of x and checks  
whether message contain signature  
or not.  
if the msg contain signature he  
will encrypt the msg with his private  
key and send it to y. then y accept  
the msg based on trust.