

Tautology:-

is a formula or assertion that is true in every possible interpretation for example $x = y$ or $x \neq y$

$p: x = y$	$q: x \neq y$	$p \vee q$
T	F	T
F	T	T

Tautology

i.e.; tautology is a statement that is always true

Contradiction:-

is a statement that is always false

p	q	$p \wedge q$
T	F	F
F	T	F

logical equivalence:

the two statements are logically equivalent if they have the same truth table. for example if $p \equiv p$ and $q \equiv \sim(\sim p)$

p	$\sim p$	$\sim(\sim p)$
T	F	T
F	T	F

that is, $p \equiv q$

logical conclusion:-

let P and Q are phrases, we say that P be a logically
requiring from Q or Q implies logically from P and
written by $P \Rightarrow Q$ if $P \rightarrow Q$ tautology and if Q isn't
conclude from P and written by $P \not\Rightarrow Q$

Ex:-

$$\text{let } P: (P \rightarrow Q) \wedge (Q \rightarrow R)$$

$$Q: P \rightarrow R$$

Is $P \Rightarrow Q$?

Sol:-

P	Q	R	$P \rightarrow Q$	$Q \rightarrow R$	$(P \rightarrow Q) \wedge (Q \rightarrow R)$	$P \rightarrow R$	$P \Rightarrow Q$
T	T	T	T	T	T	T	T
T	T	F	T	F	F	F	T
T	F	T	F	T	F	T	T
T	F	F	F	T	F	F	T
F	T	T	T	T	T	T	T
F	T	F	T	F	F	T	T
F	F	T	T	T	T	T	T
F	F	F	T	T	T	T	T

$\therefore P \Rightarrow Q$ because $P \rightarrow Q$ is a tautology

Definitions:-

let P and Q be phrases then:

- ① $Q \rightarrow P$ be the converse of $P \rightarrow Q$
- ② $\sim Q \rightarrow \sim P$ be the contrapositive of $P \rightarrow Q$.

Theorem:-

let P and Q be two statements

- ① $P \Rightarrow Q$ iff $\sim P \vee Q$ is a tautology.
- ② $P \Rightarrow Q \wedge Q \Rightarrow P$ then $P \equiv Q$.

proof:-

① let $P \Rightarrow Q$ we have to prove that $\sim P \vee Q$ be a tautology

Since $P \Rightarrow Q$ then $P \rightarrow Q$ is a tautology

then $P \rightarrow Q \equiv \sim P \vee Q$ (i.e.)

then $\sim P \vee Q$ is a tautology

Conversely:-

let $\sim P \vee Q$ be a tautology we have to prove that $P \Rightarrow Q$

(i.e. $P \rightarrow Q$ is a tautology)

Since $\sim P \vee Q$ is a tautology

then $P \rightarrow Q \equiv \sim P \vee Q$ and since $\sim P \vee Q$ is a tautol-

ogy then $P \rightarrow Q$ is a tautology and $P \Rightarrow Q$

Ex:-

Prove that $p \Rightarrow q \text{ iff } \sim q \Rightarrow \sim p$

Sol:-

Since $p \Rightarrow q$ then $p \rightarrow q$ is a tautology

Since $\sim q \rightarrow \sim p \equiv p \rightarrow q$

Then $\sim q \rightarrow \sim p$ is a tautology that is $\sim q \Rightarrow \sim p$

Conversly:-

Let $\sim q \Rightarrow \sim p$

Since $\sim q \Rightarrow \sim p$ is a tautology then $\sim q \rightarrow \sim p$ is a tautology

$\therefore p \rightarrow q \equiv \sim q \rightarrow \sim p$ that is $p \rightarrow q$ is a tautology

thus $p \Rightarrow q$

Algebra of Statements

Suppose that p and q be statements then:

1- Identity property

$$\textcircled{a} p \wedge q \equiv p \quad \textcircled{b} p \vee p \equiv p$$

2- Commutativity property

$$\textcircled{a} p \vee q \equiv q \vee p \quad \textcircled{b} p \wedge q \equiv q \wedge p$$

3- Associativity property

$$\textcircled{a} (p \vee q) \vee r \equiv p \vee (q \vee r)$$

$$\textcircled{b} (p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$$

4- Distributivity property :-

$$a) p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

$$b) p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

5- Identity property :-

$$a) p \vee I \equiv I \quad b) p \vee 0 \equiv p \quad \text{s.t } I \text{ is a tautology and } 0 \text{ is a contradiction}$$

6- Complementarity property :-

$$a) p \vee \sim p \equiv I \quad b) p \wedge \sim p \equiv 0 \quad c) \sim(\sim p) \equiv p$$

$$d) \sim I \equiv 0 \quad e) \sim 0 \equiv I$$

s.t I is a tautology statement and 0 is a contradiction statement.

7- Demorgan laws :-

$$a) \sim(p \wedge q) \equiv \sim p \vee \sim q$$

$$b) \sim(p \vee q) \equiv \sim p \wedge \sim q$$

Remarks :-

1) we can prove any of the above properties using truth table

$$2) (p \rightarrow q) \equiv (\sim p \vee q) \quad \text{(an established fact)}$$

Ex: -

prove without using truth table the following statement

$$\sim(p \rightarrow q) \equiv p \wedge \sim q$$

proof: -

$$\text{since } p \rightarrow q \equiv \sim p \vee q \text{ then } \sim(p \rightarrow q) \equiv \sim(\sim p \vee q)$$

$$\sim(p \rightarrow q) \equiv p \wedge \sim q$$

Existential quantifier and universal quantifiers

Existential quantifiers:-

Suppose that $p(x)$ is open sentence in X over the set A then the statement 'there exist $x \in A$ s.t. $p(x)$ is true' is called existential quantifier and denoted by

$\exists x \in A \ni p(x)$ is true and the symbol \exists is there exist and \ni is (s.t.) such that:

Universal property

Suppose that $p(x)$ is an open set in X over A then the statement 'for all $x \in A$ s.t. $p(x)$ is true' is called universal quantifiers and denoted by $\forall x \in A \ni p(x)$ is true and the symbol \forall is for all.

Ex:-

$(\forall x \in \mathbb{R} \ni x^2 \geq 0)$ is a universal quantifiers and true

Ex:

There exist $x \in \mathbb{N}$ s.t. $x = 7$

∴ existential quantifier and true

Ex:-

$\exists x \in \mathbb{N} : x + 2 = 7$

existential quantifier and true

Ex:-

$\exists x \in \mathbb{N} : x^2 = 1 \wedge x > 4$

existential quantifier and false

(1) Negation of statement which containing of universal quantifiers

$$(\forall x \in A : p(x)) \equiv \exists x \in A : \neg p(x)$$

(2) Negation of statement which containing of existential quantifier

$$(\exists x \in A : p(x)) \equiv \forall x \in A : \neg p(x).$$

Ex:-

What is the negation of the following statement:

$$\forall x \in \mathbb{R} : (x^2 - 1 > 0) \vee (x^2 + x > 0 \rightarrow x > 0)$$

$$\exists x \in \mathbb{R} : (x^2 - 1 \leq 0) \wedge (x^2 + x \leq 0 \nrightarrow x \leq 0)$$

Def:-

Argument: let S_1, S_2, \dots, S_n be a statement and S be a statement conclu from the statement S_1, S_2, \dots, S_n then S be a argument and denoted by $S_1, S_2, \dots, S_n \vdash S$

Remark:- The argument $S_1, S_2, \dots, S_n \vdash S$ be true iff $S_1 \wedge S_2 \wedge S_3 \dots \wedge S_n \rightarrow S$ is a tautology

$$\text{Q} p \vee \neg q, p, \neg q \vdash p \vee q$$

Sol:-

$$S_1 = p \vee q, S_2 = \neg q, S = p \vee q$$

p	q	$\neg p$	$\neg q$	$p \vee \neg q$	$p \vee q$	$S_1 \wedge S_2$	$S_1 \wedge S_2 \rightarrow S$
T	T	F	F	T	T	F	F
T	F	F	T	T	T	T	T
F	T	T	F	F	F	F	F
F	F	T	T	T	F	T	T